Artificial Intelligence V04: Local and adversarial search



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From hill climbing search to genetic algorithms Game playing Resource limits and other difficulties

Based on material by Stuart Russell, UC Berkeley





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2048 leaderboard link



https://goo.gl/meh3Ro

Educational objectives

- Re-tell the story of improving local search from hill-climbing to genetic algorithms
- **Remember** the **minimax**, α - β and **expectiminimax** algorithms
- Implement an Al agent for a given simple game

"In which we relax the simplifying assumptions of the previous lecture, thereby getting closer to the real world; including the problems that arise when we try to plan ahead in a world where other agents are planning against us."

→ Reading: AIMA, [ch. 4.1-4.2 (local search)]; ch. 5 (games)





1. FROM HILL CLIMBING SEARCH TO GENETIC ALGORITHMS

Local search

Example: *n*-queens problem



Task

• Put *n* queens on a $n \times n$ board with no two queens on the same row, column, or diagonal



Possible solution

- Initialize one queen per column
- Move one queen up/down at a time to reduce number of conflicts using heuristic h
- Almost always solves n-queens problems almost instantaneously (#states: n^n)
 - → works for very large n, e.g., n = 1'000'000

Iterative improvement algorithms



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Local search: search for optimal states instead of path's

- In many optimization problems, path is irrelevant; the goal state itself is the solution
 - → State space: set of "complete" configurations;
 - → Goal: find optimal configuration (or a configuration satisfying constraints)
- Examples: TSP, timetable

Iterative improvement

- In such cases: use iterative improvement algorithms
 - → Keep a single "current" state, try to improve it
 - → Constant space, suitable for online as well as offline search

Possible implementations

- Hill climbing
- Simulated annealing
- Genetic algorithms

Hill climbing search (a.k.a. gradient ascent/descent)



Systematic search for an optimum

- Analogy: «Like climbing Everest in thick fog with amnesia»
- Result: finds a state that is a local maximum

...by selecting only the highest-valued successor for expansion iif its value is better

The state space landscape

- Practical problems typically have an **exponential number of local maxima** to get stuck in
- Random-restart hill climbing overcomes local maxima → trivially complete
- Random sideways moves escape from shoulders (good), loop on flat maxima (bad)





Hill climbing search: an outlook

All previously discussed search algorithms only work in discrete state and action spaces (otherwise the branching factor is infinite)

- Hill climbing in continuous space (gradient descent) is the work horse of deep learning ٠
- Some pointers: •
 - → http://sebastianruder.com/optimizing-gradient-descent/ (overview of the gradient descent family)
 - → https://stdm.github.io/Some-places-to-start-learning-ai-ml/ (links to courses on deep neural networks)

PROC. OF THE IEEE, NOVEMBER 1998		
Gradient-Based Learnin	g Applied to Document	
Recog	nition	
Yann LeCun, Léon Bottou, Yoshua Bengio, and Patrick Haffner		
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HMM Hidden Markov model. HOS Heuristic oversegmentation.	method of recognizing individual patterns consists in divid- ing the system into two main modules shown in figure 1. The first module, called the feature extractor, transforms	
K-NN K-nearest neighbor. NN Neural network. Order lobum atom measuritien	the input patterns so that they can be represented by low- dimensional vectors or short strings of symbols that (a) can	
 OUR Optical character recognition. PCA Principal component analysis. RBF Radial basis function. 	be easily matched or compared, and (b) are relatively in- variant with respect to transformations and distortions of the invariant mattemps that do not shown their nature. The	
BS-SVM Reduced-set support vector method. SDNN Space displacement neural network. SVM Support vector method. TDNN Time delay neural network. V-SVM Virtual support vector method.	the input patterns that do not change their nature. The feature extractor contains most of the prior knowledge and is rather specific to the task. It is also the focus of most of the design effort, because it is often entirely hand-trafted. The classifier, on the other hand, is often general-putpose	

uned by the ability of the designer to come up with a appropriate set of features. This turns out to be a daux task which, unfortunately, must be redone for each n



Dearn		i statistication of the second
	D;	y gradient descent
Marcia Andrychowicz ¹ , Misha Denil ¹ , Sergio Gómez Colmenarejo ¹ , Matthew W, Hoffman ¹ , David Plau ¹ , Tan Schau ¹ , Brendan Shillineford ^{1,2} , Nanda de Freilar ^{1,2,3}		
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	Abstr	act
been wildly succ by hand. In this p cast as a learning the problems of i by LSTMs, output they are trained, demonstrate this neural networks,	essful. In spite of this, op apper we show how the des problem, allowing the al netrest in an automatic wa erform generic, hand-desij and also generalize well i on a number of tasks, incl and styling images with r	timization algorithms are still designed ign of an optimization algorithm can be lgorithm to learn to exploit structure in y. Our learned algorithms, implemented gand competitors on the tasks for which to new tasks with similar structure. We usling simple convex problems, training neural art.
Introduction		
Frequently, tasks in mach function $f(\theta)$ defined ov $\beta^* = \arg \min_{\theta \in \Theta} f(\theta)$. I applied, the standard appro- n a sequence of updates	ine learning can be expre er some domain $\theta \in \Theta$. While any method capabl each for differentiable fun	ssed as the problem of optimizing an objective The goal in this case is to find the minimizer de of minimizing this objective function can be ctions is some form of gradient descent, resulting
	$\theta_{t+1} = \theta_t - c$	$s_t \nabla f(\theta_t)$.
he performance of vanill f gradients and ignores s ehavior by rescaling the f second-order partial de tatrix or Fisher informati	a gradient descent, howev second-order information gradient step using curvate rivatives—although other on matrix are possible.	er, is hampered by the fact that it only makes use 0. Classical optimization techniques correct this ure information, typically via the Hessian matrix choices such as the generalized Gauss-Newton

2016

30.7

arXiv:1606.04474v2 [cs.NE]

98], Rprop [Riedminer and Linken, 2012], and ADAM [K

on on Named Information Descarcing Systems (MIDE 2016), Researchers, St.

Simulated annealing Towards optimizing hill climbing search

Idea (by [Metropolis et al., 1953] for physical process modelling)

- Escape local maxima by allowing some "bad" moves
- ...but gradually decrease their size and frequency

Application

- For "good" schedule of decreasing the temperature (→ see appendix), it always reaches the best state
- Widely applied for e.g. VLSI layout, airline scheduling

Modern variants

 «momentum», «Adam» and other adaptation strategies for a «learning rate»
 → first link on the last slide





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Local beam search

...and still optimizing hill climbing search

Idea

- Keep k states instead of 1: choose top k of all their successors • (not the same as k searches run in parallel! \rightarrow Why?)
- Searches that find good states recruit other searches to join them •

Problem

Quite often, all k states end up on same local hill

Idea contd

- **Choose** k successors randomly, biased towards good ones
 - → Observe the close analogy to natural selection!





Not: each



Compare Don Knuth on "the advantages of unbiased sampling as a way to gain insight into a complicated subject" (e.g. ch. 2 in the above book)



Genetic algorithms (GA)

... improving on the idea of local beam search



Idea

- Combine stochastic local beam search + generating successors from *pairs* of states
 - → uphill tendency + random exploration + exchange of information among searches



Example: 8-queens states encoded as digit strings. The original population (left) is ranked by a fitness function, resulting in pairs for mating. The offspring is subject to mutation.

Application

- GAs require states encoded as strings
- Crossover helps iif substrings are meaningful components
- GAs \neq evolution



2. GAME PLAYING

Adversarial search

Games vs. search problems

"Unpredictable" opponent

→ solution is a strategy (specifying a move for every possible opponent reply)

Time limits

→ unlikely to find goal, must approximate



- Computer considers possible lines of play (computer chess: Babbage, 1846)
- Algorithm for perfect play (minimax: Zermelo, 1912; game theory: von Neumann, 1944)
- Finite horizon, approximate evaluation (depth cut-off: Zuse, 1945; Wiener, 1948; Shannon, 1950)
- First chess program (Turing, 1951)
- Machine learning to improve evaluation accuracy (RL for checkers: Samuel, 1952-57)
- Pruning to allow deeper search (α - β search: McCarthy, 1956)







Types of games



	deterministic	stochastic
perfect information	chess , checkers («Dame»), go, othello («Reversi»)	backgammon, monopoly
only partial observability	battleship , kriegspiel (chess without seeing enemy pieces)	bridge (~ «Jass», «Skat»), poker , scrabble, <i>global</i> <i>thermonuclear war</i>



Deterministic (turn-based, 2-player) games The search tree, e.g. of tic-tac-toe

Utility

-1

0

+1

"Max"

Plaver's name

- 1st player (moves first)
- Wants to maximize utility of terminal states
- Tree shows Max's perspective

"Min"

2nd player
Wants to minimize (Max's) utility

Utility

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- Numeric value ("payoff") of terminal state
- "zero-sum game" iif total payoff (to all players) is constant over all game instances





Minimax: depth-first exploration of game tree Perfect play for deterministic, fully observable games

Idea

- Choose move to position with highest minimax value
- Minimax value: highest value among options minimized by adversary
 - → best achievable payoff against best play

Example

- Any 2-ply game tree (i.e., each player moves once)
- Max's best move at root: a₁ (leading to highest minimax value of 3)
- ...because Min's best reply will be b₁ (leading to lowest minimax value / utility of 3)





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Minimax (contd.) Algorithm and properties

function Minimax-Decision(state) returns an action
inputs: state, current state in game
return the a in Actions(state) maximizing Min-Value(Result(a, state))

function Max-Value(state) returns a utility value
if Terminal-Test(state) then return Utility(state)
v ← -∞
for a, s in Successors(state) do v ← Max(v, Min-Value(s))
return v

function Min-Value(state) returns a utility value
if Terminal-Test(state) then return Utility(state)
v ← ∞
for a, s in Successors(state) do v ← Min(v, Max-Value(s))
return v



Properties

- Complete? Yes, if tree is finite (e.g., chess has specific rules for this)
- Optimal? Yes, against an optimal opponent (Otherwise?)
- Time complexity? $O(b^m)$
- Space complexity? O(bm) (for depth-first exploration)
- For chess, b = 35, m = 100 for "reasonable" games
 - → exact solution completely infeasible; but do we need to explore every path?

α - β pruning example Overcoming exponential (b^m) number of states to be explored



Successively tightening bounds on minimax values

- *α* is the best value (to Max) found so far in current subtree of a Max node (A→α)
- If any node v is worse than α, Max will not choose it
 > prune that **branch**
- Similarly: β is best score Min is assured of in current subtree of a Min node (B, C, D $\rightarrow \beta$)





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α - β pruning example Overcoming exponential (b^m) number of states to be explored



[lower bound, uper bound]

_∞, +∞] A

B has at most a

value of 3

(a)

Recursion: B
$\alpha = -\infty$
$\beta = 3$

Successively tightening bounds on minimax values

- α is the best value (to Max) found so far in current subtree of a Max node (A $\rightarrow \alpha$)
- If any node v is worse than α, Max will not choose it
 → prune that branch
- Similarly: β is best score Min is assured of in current subtree of a Min node (B, C, D $\rightarrow \beta$)



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	Recursion: A
	$\alpha = 3$
l	$\beta = \infty$



Zurich University of Applied Sciences and Arts InIT Institute of Applied Information Technology (stdm) Successively tightening bounds on minimax values

- *α* is the best value (to Max) found so far in current subtree of a Max node (A→α)
- If any node v is worse than α, Max will not choose it
 → prune that **branch**
- Similarly: β is best score Min is assured of in current subtree of a Min node (B, C, D→ β)

α - β pruning example Overcoming exponential (b^m) number of states to be explored

Successively tightening bounds on minimax values

- α is the **best value** (to Max) found so far in current subtree of a Max node (A $\rightarrow \alpha$)
- If any node v is worse than α , Max will not choose it **→** prune that branch
- Similarly: β is best score Min is assured of in current subtree of a Min node (B, C, D $\rightarrow \beta$)



[lower bound, uper bound]

[_∞, +∞] 🛦

B has at most a

value of 3

B has

exactly

value 3, so A has

at least 3

(a)

[-∞, 3]

3

[3, 3]

3 12

B



Recursion: C

 $\alpha = 3$ $\beta = 2$



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[3, 3] 🗑

12 3

α - β pruning example Overcoming exponential (b^m) number of states to be explored

[lower bound, uper bound]



[<u>-∞, +∞]</u> A

B has at most a

value of

exactly

value 3. so A has

at least 3

(a)



Recursion: D

 $\alpha = 3$ $\beta = 14$



Successively tightening bounds on minimax values

- α is the best value (to Max) found so far in current subtree of a Max node $(A \rightarrow \alpha)$
- If any node v is worse than α , Max will not choose it → prune that branch
- Similarly: β is best score Min is assured of in current subtree of a Min node (B, C, D $\rightarrow \beta$)

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α - β pruning example Overcoming exponential (b^m) number of states to be explored

[3, 3]

3

(d)



→ prune that branch

[3, +∞] ▲

[-∞, 2]

2



[lower bound, uper bound]

[<u>-∞, +∞]</u> A

B has at most a

value of 3

(a)

(b)

[-∞, 3]

3

[-∞, 3]

3 12

B

Similarly: β is best score Min is assured of in current subtree of a Min node (B, C, D $\rightarrow \beta$)

Successively tightening bounds on minimax values

- α is the best value (to Max) found so far in current subtree of a Max node $(A \rightarrow \alpha)$

If any node v is worse than α , Max will not choose it



2



[2, 2]

14

5

2

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Algorithm and properties (changes to minimax in bold-italic)



 α - β pruning (contd.)

function Alpha-Beta-Search(state) returns an action

```
function Min-Value(state, \alpha, \beta) returns a utility value
     if Terminal-Test(state) then return Utility(state)
     v \leftarrow \infty
     for a in Actions(state) do
          v \leftarrow Min(v, Max-Value(Result(state, a), \alpha, \beta))
          if v < \alpha then return v
          \beta \in Min(\beta, v)
     return v
```

Properties

return v

- Pruning does not affect final result
- Good move ordering improves effectiveness of pruning ٠
- With "perfect ordering", time complexity = $O(b^{m/2})$ ٠
 - → doubles solvable depth
 - → a simple example of the value of reasoning about which computations are relevant (metareasoning)
 - \rightarrow unfortunately, $35^{100/2}$ (for chess) is still impossible!





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3. RESOURCE LIMITS AND OTHER DIFFICULTIES

Adversarial search (contd.)

Resource limits Towards real-world conditions

Standard approach

- Use **Cutoff-Test** instead of Terminal-Test e.g., depth limit (perhaps add **quiescence search**: only cut off search at positions that don't drastically change their value in the near future, e.g. captures in chess; otherwise continue search)
- Use **Eval** instead of Utility i.e., evaluation function that estimates desirability of position
- Lookup of start/end games

Example

- Suppose we have 100 seconds, explore 10⁴ nodes/second
 → 10⁶ nodes per move ≈ 35^{8/2}
- α - β reaches depth 8
 - ➔ pretty good chess program





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Eval(uation) functions Designing or learning effective cutoff tests





Example: The two chess positions differ only in the position of the rook ("Turm") at lower right.
In (a), Black has an advantage of a knight ("Springer") and two pawns ("Bauern") → should be enough to win.
In (b), White will capture the queen ("Dame") → should be strong enough to win.

For chess, typically linear weighted sum of features

- Eval(s) = $w_1 f_1(s) + w_2 f_2(s) + ... + w_n f_n(s)$
- Example:

 $w_1=9$ and $f_1(s) = (number of white queens) - (number of black queens)$

• Can be learned with machine learning techniques

Nondeterministic (stochastic) games

Chance is introduced by e.g. dice-rolling or card-shuffling

Simplified example

- A game with coin-flipping ٠
- Nondeterminism is handled by an ٠ additional level in the tree, consisting of chance nodes



Real-world example

- 2048: numbers appear with probability $P(2) = \frac{9}{10}$ and $P(4) = \frac{1}{10}$ at random free board positions
- Backgammon: Before each move, dice-• rolls determine the legal moves



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2

4

Expectiminimax – maximizing expected value Algorithm and properties

```
function ExpectiMinimax-Decision(state) returns an action
inputs: state, current state in game
return a in Actions(state) maximizing
ExpectiMinimax-Value(Result(a, state))
```

function ExpectiMinimax-Value(state) returns a utility value

- if Terminal-Test(state) then
 return Utility(state)
- if state is a Max node then
 return highest ExpectiMinimax-Value of Successors(state)
- if state is a Min node then
 return lowest ExpectiMinimax-Value of Successors(state)
- if state is a chance node then
 return average of ExpectiMinimax-Value of Successors(state)

$$\begin{split} & \text{Expectiminimax}(s) = \\ & \begin{cases} \text{Utility}(s) & \text{if } \text{Terminal-Test}(s) \\ & \max_a \text{Expectiminimax}(\text{Result}(s, a)) & \text{if } \text{Player}(s) = \text{max} \\ & \min_a \text{Expectiminimax}(\text{Result}(s, a)) & \text{if } \text{Player}(s) = \text{min} \\ & \sum_x P(r) \text{Expectiminimax}(\text{Result}(s, r)) & \text{if } \text{Player}(s) = \text{chance} \end{cases} \end{split}$$

Properties

- Algorithm works just like Minimax except chance-nodes are also handled
- Expectiminimax gives perfect play
- In case of only 1 player, Expectiminimax becomes Expectimax
- Time complexity: $O(b^m n^m)$ (where n is the number of distinct random events, e.g. dice rolls)
 - → Possibilities are multiplied enormously in games of chance
 - → Simultaneously, no likely sequences exist to do effective α - β pruning



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Nondeterministic games in practice

Example Backgammon

- Dice rolls increase *b* (21 possible rolls with 2 dice)
- Ca. 20 legal moves (can be 6,000 with 1-1 roll)
 → at depth 4: 20 × (21 × 20)³ ≈ 1.2 × 10⁹ nodes
- As depth increases, probability of reaching a given node shrinks
 - → value of lookahead is diminished (see next slide for consequences)

But

- «TDGammon» (Tesauro, 1992) uses depth-2 search \approx world-champion level
 - → uses neural network and reinforcement learning (RL) to train Eval function via games against itself

predicted probability



• → see also <u>http://webdocs.cs.ualberta.ca/~sutton/book/ebook/node108.html</u>



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Digression: Exact Eval values do matter ...for nondeterministic games





The vastly enlarged scale (compared to the left tree) of the utility value suffices to form the highest expectiminimax value despite the low probability

- Behavior is preserved only by positive linear transformation of Eval
- Hence Eval should be proportional to the expected payoff

→ Exact values don't matter for deterministic games → see appendix

Games of imperfect information

E.g., card games, where opponent's initial cards are unknown

First idea («averaging over clairvoyance»)

- Typically we can calculate a probability for each possible deal
 - → Seems just like having one big dice roll at the beginning of the game *
 - → Idea: compute the Minimax value of each action in each deal;
 - then, choose the action with highest expected value over all deals (i.e., Expectiminimax) *

Monte Carlo simulation alone does not always suffice: it handles randomness well (e.g. in the deal), but not strategy that looks like randomness (e.g. an adversary's move in Kriegspiel)

«GIB» - Ginsberg's Intelligent Bridgeplayer

- Current best bridge program (Ginsberg, 1999)
- Approximates idea above by Monte Carlo simulation of some of the possible deals to be computationally feasible
 - 1. Generate 100 deals consistent with bidding information
 - 2. Pick the action that wins most tricks on average

Skat: a wonderful partially observable stochastic card game with some similarity to Bridge → see <u>https://www.pagat.com/schafk/skat.html</u>

iminimax) *



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NAJA, MACHT NIX, ER HAT SOWNESD VERLOREN

*) Proper analysis of averaging over clairvoyance

Intuition might mislead

- That the value of an action is the average of its values in all possible states is WRONG
 (→ see appendix)
- With partial observability, value of an action depends on the belief state the agent will end up in

Maintaining its belief state is a **core function of any intelligent system** in partially observable (~real world) environments → see AIMA, ch. 4.4 The agent's **current information about** all the **possible physical states** it might be in (given the sequence of actions and percepts up to that point)

Dealing with belief states

- Generate and search a tree of belief states
- · Leads to rational behaviors such as
 - Acting to obtain information
 - Signaling to one's partner
 - Acting randomly to minimize information disclosure



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Where's the intelligence?

- Searching over the right state spaces leads to **emergent "clever" behavior** in games (information gathering, alliances, bluffs)
- Many local and adversarial search methods can be **enhanced by learning** (e.g., learn good Eval functions by playing games against oneself)
- The brain might perform local (hill climbing) search as well
 - → see https://www.cs.toronto.edu/~hinton/backpropincortex2014.pdf

→ Hill-climbing, minimax, α - β , etc. are still pure computation → Intelligent behavior emerges from their composition on **suitable** data structures







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Review

- Local search algorithms evolve a small number of states towards better utility (typical: 1 state)
- In **continuous space**, local search by linear programming or convex optimization is extremely efficient in practice (polynomial time complexity!)
- In non-deterministic environments, keeping track of one's belief state is paramount
- Games are fun to work on and dangerous
- They illustrate several important points about AI
 - perfection is unattainable → must approximate
 - good idea to **think about what to think** about (metareasoning → pruning)
 - uncertainty constrains the assignment of values to states
 - optimal decisions depend on information state (belief state), not real state
- → Games are to AI as grand prix racing is to automobile design







APPENDIX

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Pseudocode for hill climbing search and simulated annealing



```
function Hill-Climbing (problem) returns a state that is a local maximum
   inputs: problem, a problem
   local variables: current, a node
                    neighbour, a node
   loop do
       neighbour \leftarrow a highest-valued successor of current
       if Value[neighbour] < Value[current] then return State[current]
       current ← neighbour
   end
function Simulated-Annealing (problem, schedule) returns a solution state
   inputs: problem, a problem
           schedule, a mapping from time to "temperature"
   local variables: current, a node
                    next, a node
                    T, a "temperature" controlling the probability of downward steps
   for t \leftarrow 1 to \infty do
       T \leftarrow schedule[t]
       if T = 0 then return current
       next \leftarrow a randomly selected successor of current
       ΔE ← Value[next]-Value[current]
       if \Delta E > 0 then current \leftarrow next
```

```
else current \leftarrow next only with probability e^{\frac{1}{T}}
```

Digression: Exact Eval values don't matter ...for deterministic games





- Behavior is preserved under any **monotonic** transformation of Eval
- Only the order matters: payoff in deterministic games acts as an ordinal utility function

Deterministic games in practice

Checkers

- *«Chinook»* ended 40-year-reign of human world champion Marion Tinsley in 1994
- Used an endgame database defining perfect play for all positions involving 8 or fewer pieces on the board (total: 443,748,401,247)

Chess

- «Deep Blue» defeated human world champion Gary Kasparov in a six-game match in 1997
- Searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply

Othello

• Human champions refuse to compete against computers, who are too good

Go

- 2010: human champions refuse to compete against computers, who are too bad
- 2016: Google DeepMind's «AlphaGo» unexpectedly defeats world champion Lee Sedol
- $b > 300 \rightarrow$ most programs use pattern knowledge bases to suggest plausible moves











Expectations & expected values adapted from U Washington`s CSE473, lecture 8



The expectation operator E()

- We can define a function f(X) of a random variable X
- The expected value of a function is its average value under the probability distribution over the function's inputs:

$$E(f(X)) = \sum_{x} f(X = x)P(X = x)$$

Example

- How long to drive to the airport?
- Driving time D (in mins) as a function of traffic T: D(T = none) = 20, D(T = light) = 30, D(T = heavy) = 60
- What is your expected driving time? Let probability $P(T) = \{none: 0.25, light: 0.5, heavy: 0.25\}$ $\Rightarrow E(D(T)) = D(none)P(none) + D(light)P(light) + D(heavy)P(heavy)$ $\Rightarrow E(D(T)) = 20 * 0.25 + 30 * 0.5 + 60 * 0.025 = 35 mins$

Why averaging over clairvoyance is wrong A common sense example of a journey

Day 1

- Road A leads to a small heap of gold pieces
- Road B leads to a fork:
 - the left fork leads to a bigger heap of gold;
 - take the right fork and you'll be run over by a bus.
- → "B" (and then "left") is the optimal choice

Day 2

- Road A leads to a small heap of gold pieces
- Road B leads to a fork:
 - take the left fork and you'll be run over by a bus;
 - the right fork leads to a bigger heap of gold.
- \rightarrow "B" (and then "right") is the optimal choice

Day 3

- Road A leads to a small heap of gold pieces
- Road B leads to a fork:
 - guess correctly and you'll find a bigger heap of gold;
 - guess incorrectly and you'll be run over by a bus.
- → "B" still seems optimal; but this ignores the resulting belief state (that includes ignorance & possibility of death!)



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