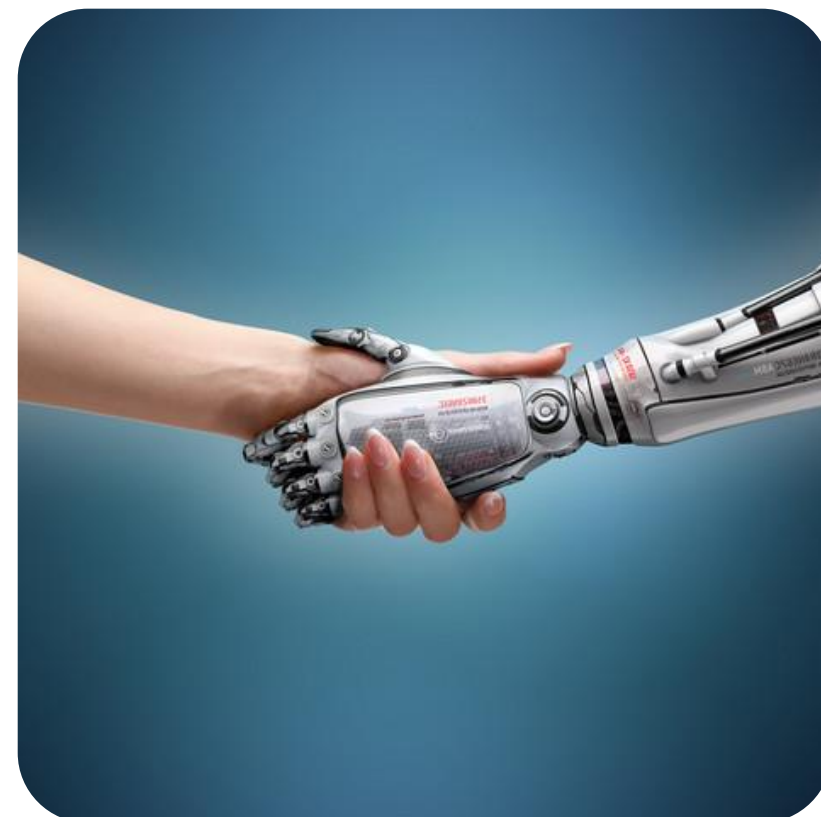


# Artificial Intelligence

## V04: Local and adversarial search

From hill climbing search to genetic algorithms  
Game playing  
Resource limits and other difficulties



Based on material by Stuart Russell, UC Berkeley



**2048 leaderboard link**

**<https://goo.gl/meh3Ro>**

# Educational objectives

- **Re-tell** the story of improving **local search** from hill-climbing to genetic algorithms
- **Remember** the **minimax**,  $\alpha$ - $\beta$  and **expectiminimax** algorithms
- **Implement** an **AI agent** for a given simple game

*“In which we relax the simplifying assumptions of the previous lecture, thereby getting closer to the real world; including the problems that arise when we try to plan ahead in a world where other agents are planning against us.”*

➔ Reading: AIMA, [ch. 4.1-4.2 (local search)]; ch. 5 (games)



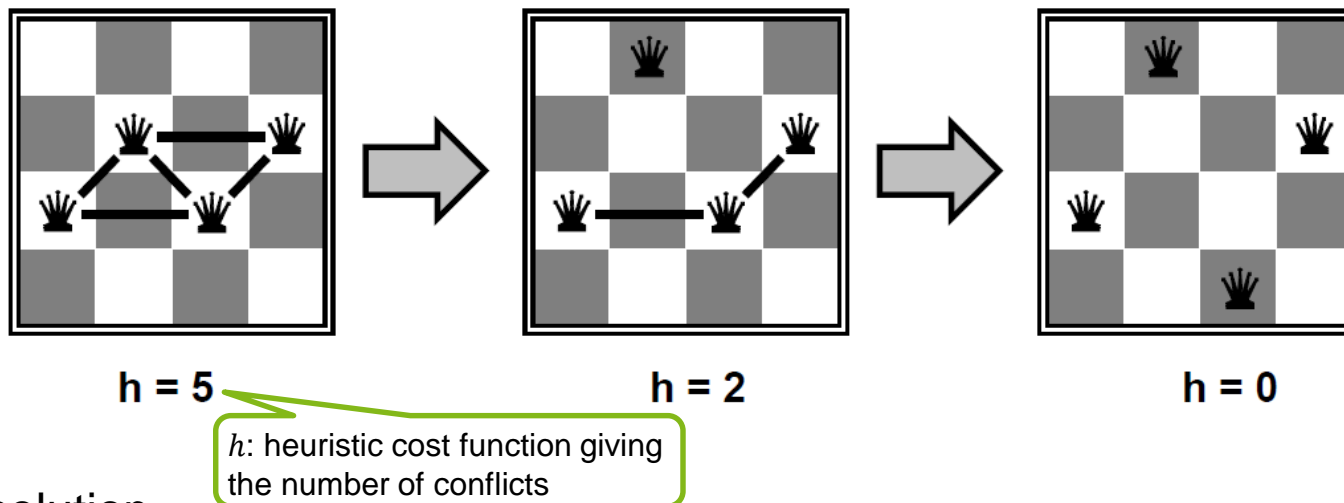
# 1. FROM HILL CLIMBING SEARCH TO GENETIC ALGORITHMS

## Local search

# Example: $n$ -queens problem

## Task

- Put  $n$  queens on a  $n \times n$  board with no two queens on the same row, column, or diagonal



## Possible solution

- Initialize one queen per column
- Move **one queen up/down at a time** to reduce number of conflicts using heuristic  $h$
- Almost always solves  $n$ -queens problems almost instantaneously (#states:  $n^n$ )  
→ works for very large  $n$ , e.g.,  $n = 1'000'000$

# Iterative improvement algorithms

**Local search:** search for optimal states instead of path's

- In many optimization problems, **path is irrelevant**; the **goal state itself is the solution**
  - State space: set of "complete" configurations;
  - Goal: find **optimal** configuration (or a configuration satisfying constraints)
- Examples: TSP, timetable

## Iterative improvement

- In such cases: use **iterative improvement** algorithms
  - **Keep a single "current" state**, try to improve it
  - Constant space, suitable for online as well as offline search

## Possible implementations

- **Hill climbing**
- **Simulated annealing**
- **Genetic algorithms**

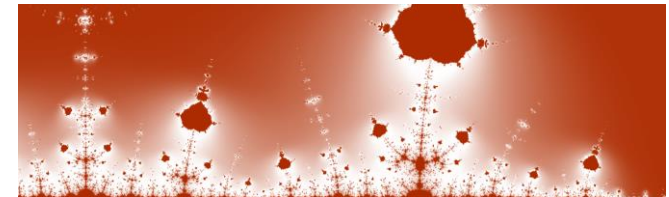
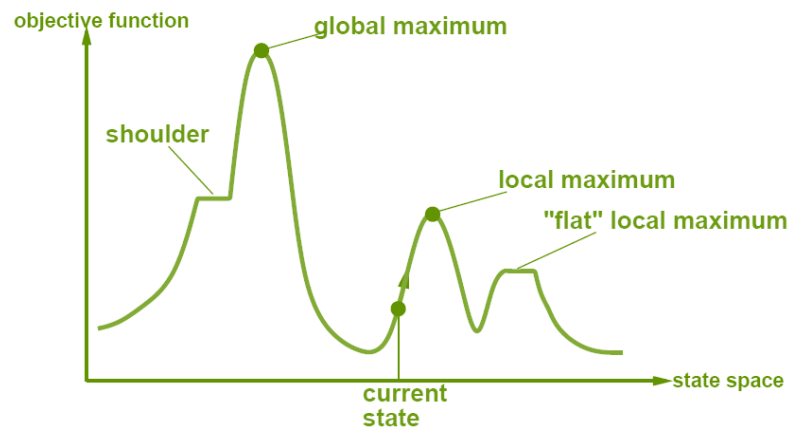
# Hill climbing search (a.k.a. gradient ascent/descent)

Systematic search for an optimum

- Analogy: «*Like climbing Everest in thick fog with amnesia*»
- Result: finds a state that is a **local maximum**  
...by selecting only the highest-valued successor for expansion **iif** its value is better

The state space landscape

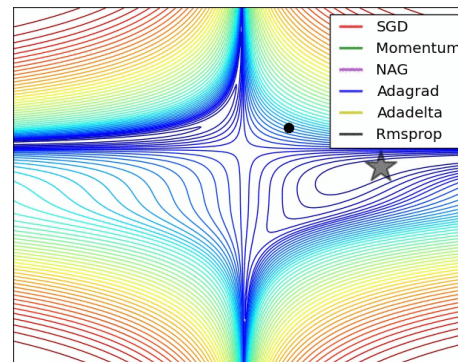
- Practical problems typically have an **exponential number of local maxima** to get stuck in
- **Random-restart** hill climbing overcomes local maxima → trivially complete
- **Random sideways** moves escape from shoulders (good), loop on flat maxima (bad)



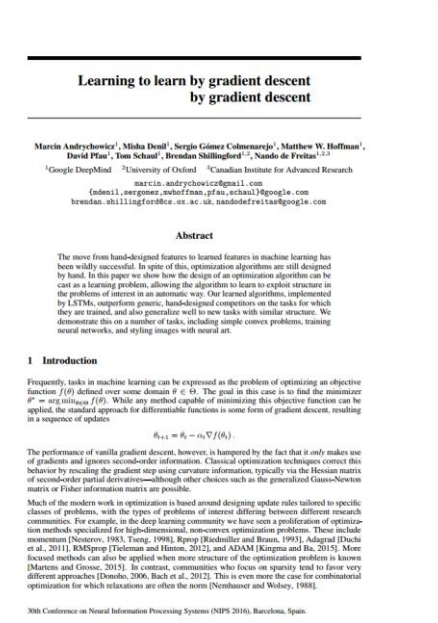
# Hill climbing search: an outlook

All previously discussed search algorithms only work in discrete state and action spaces (otherwise the branching factor is infinite)

- Hill climbing in continuous space (gradient descent) is the work horse of deep learning
- Some pointers:
  - <http://sebastianruder.com/optimizing-gradient-descent/> (overview of the gradient descent family)
  - <https://stdm.github.io/Some-places-to-start-learning-ai-ml/> (links to courses on deep neural networks)



arXiv:1606.04474v2 [cs.LG] 30 Nov 2016





# Simulated annealing

## Towards optimizing hill climbing search

Idea (by [Metropolis et al., 1953] for physical process modelling)

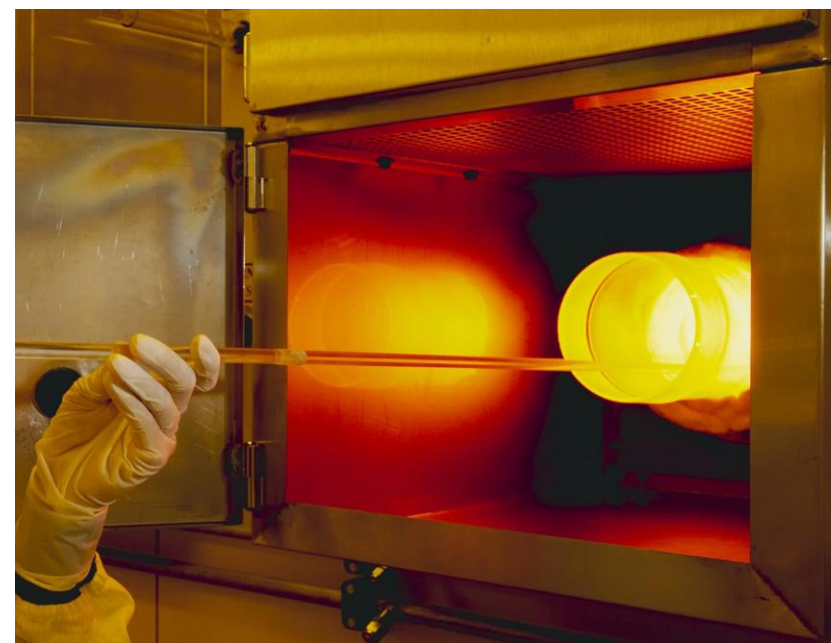
- Escape local maxima by allowing some “bad” moves
- ...but **gradually decrease** their **size and frequency**

### Application

- For “**good**” **schedule** of decreasing the temperature (→ see appendix), it **always reaches** the **best state**
- **Widely applied** for e.g. VLSI layout, airline scheduling

### Modern variants

- «**momentum**», «**Adam**» and other adaptation strategies for a «**learning rate**»  
→ first link on the last slide



# Local beam search

## ...and still optimizing hill climbing search

### Idea

- **Keep  $k$  states** instead of 1; **choose top  $k$**  of all their **successors** (not the same as  $k$  searches run in parallel! → Why?)
- **Searches** that find good states **recruit other** searches to **join** them

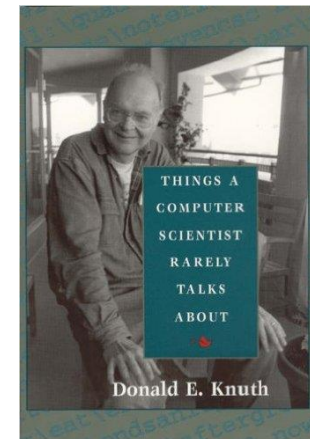
Not: each

### Problem

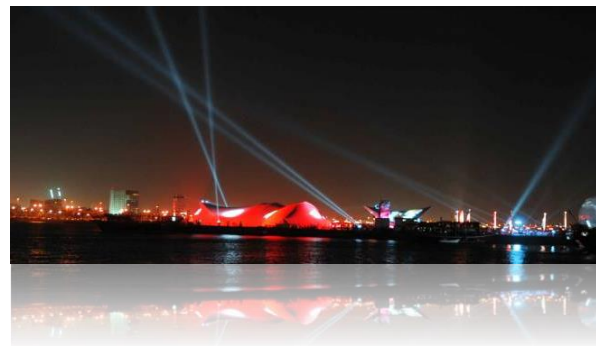
- Quite often, all  $k$  states end up on same local hill

### Idea contd.

- **Choose  $k$  successors randomly, biased towards good ones**  
→ Observe the close analogy to natural selection!



Compare Don Knuth on “the advantages of unbiased sampling as a way to gain insight into a complicated subject” (e.g. ch. 2 in the above book)

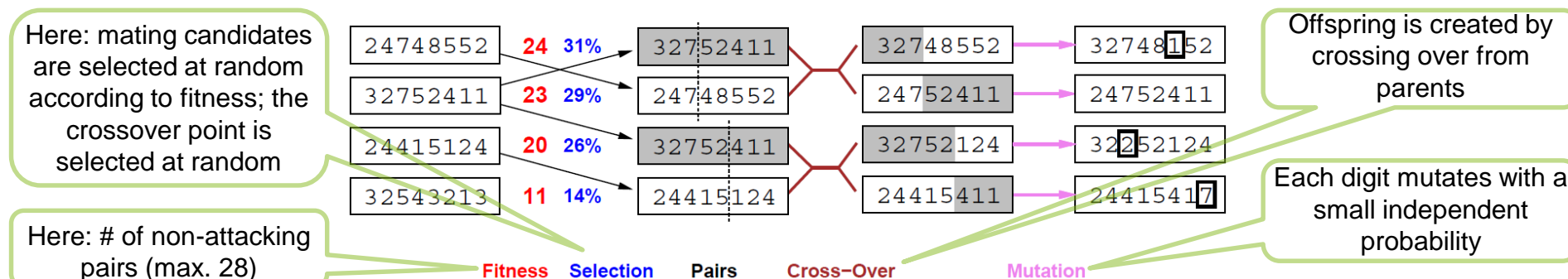


# Genetic algorithms (GA)

## ...improving on the idea of local beam search

### Idea

- Combine stochastic local beam search + generating successors from *pairs* of states  
**→ uphill tendency + random exploration + exchange of information among searches**



Example: 8-queens states encoded as digit strings. The original population (left) is ranked by a fitness function, resulting in pairs for mating. The offspring is subject to mutation.

### Application

- GAs **require states** encoded as strings
- Crossover helps **iif substrings** are meaningful components
- GAs  $\neq$  evolution

## 2. GAME PLAYING

### Adversarial search

# Games vs. search problems

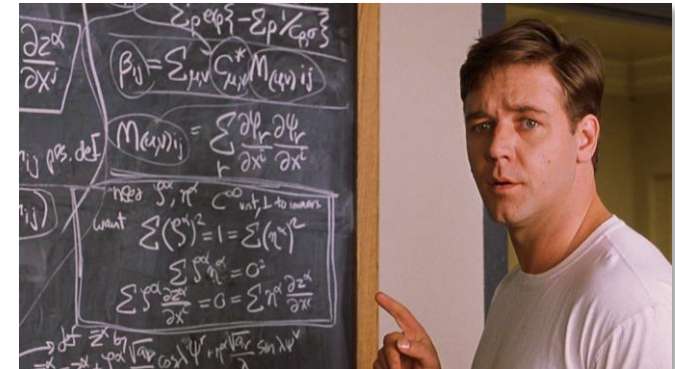
“Unpredictable” opponent

→ solution is a **strategy** (specifying a move for every possible opponent reply)

Time limits

→ unlikely to find goal, must **approximate**

Which movie?



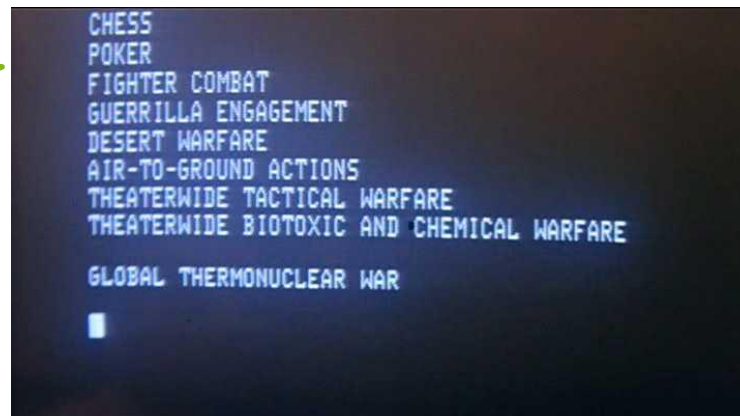
Early history:

- Computer considers possible lines of play (computer chess: Babbage, 1846)
- Algorithm for perfect play (minimax: Zermelo, 1912; game theory: von Neumann, 1944)
- Finite horizon, approximate evaluation (depth cut-off: Zuse, 1945; Wiener, 1948; Shannon, 1950)
- First chess program (Turing, 1951)
- Machine learning to improve evaluation accuracy (RL for checkers: Samuel, 1952-57)
- Pruning to allow deeper search ( $\alpha$ - $\beta$  search: McCarthy, 1956)

# Types of games

	deterministic	stochastic
perfect information	<b>chess</b> , checkers («Dame»), go, othello («Reversi»)	backgammon, <b>monopoly</b>
only partial observability	<b>battleship</b> , kriegspiel (chess without seeing enemy pieces)	bridge (~ «Jass», «Skat»), <b>poker</b> , scrabble, <i>global thermonuclear war</i>

Which movie?



# Deterministic (turn-based, 2-player) games

## The search tree, e.g. of tic-tac-toe

Player's name

„Max“

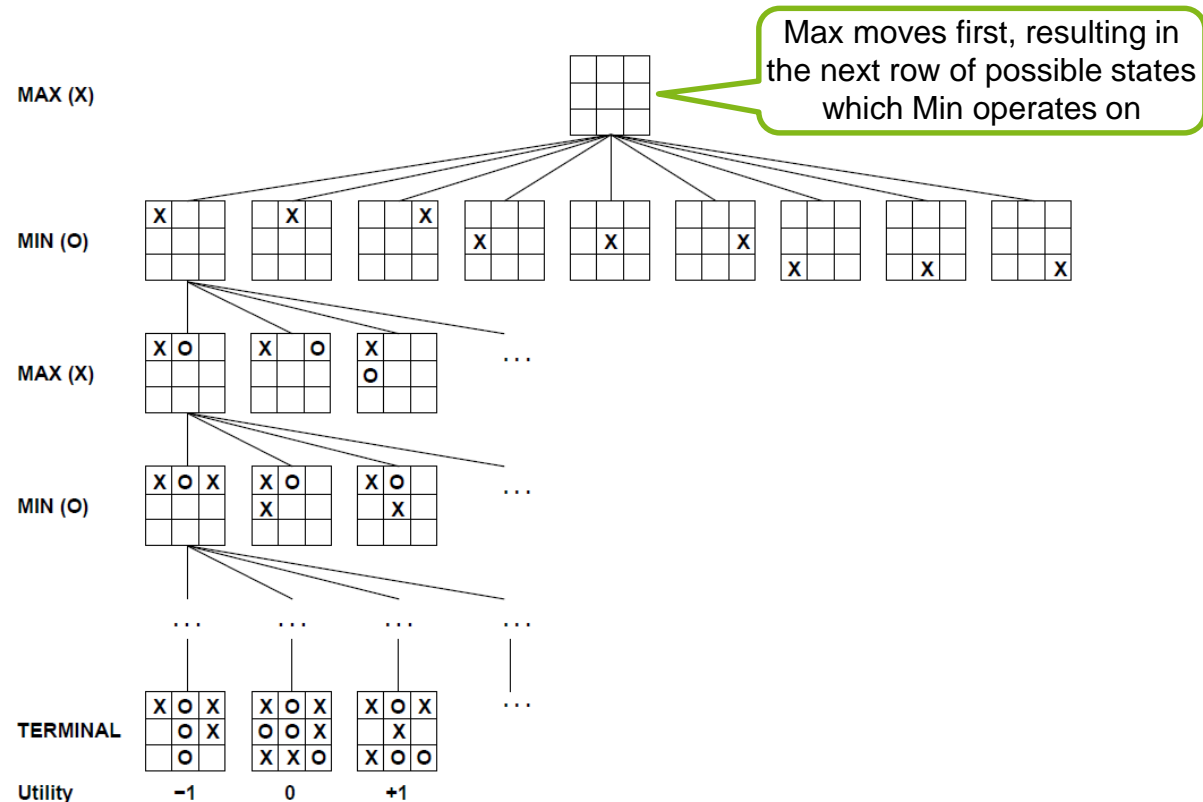
- 1<sup>st</sup> player (moves first)
- Wants to **maximize utility** of terminal states
- Tree shows Max's perspective

„Min“

- 2nd player
- Wants to **minimize (Max's) utility**

### Utility

- Numeric value („payoff“) of terminal state
- „**zero-sum game**“ iif total payoff (to all players) is constant over all game instances



# Minimax: depth-first exploration of game tree

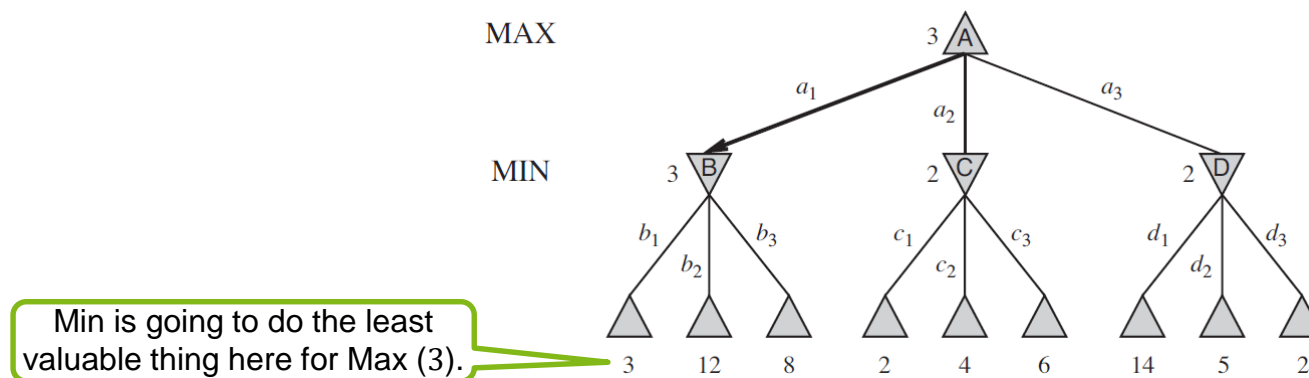
## Perfect play for deterministic, fully observable games

### Idea

- Choose move to position with highest **minimax** value
- Minimax value: **highest value among options minimized** by adversary  
→ **best** achievable payoff **against best play**

### Example

- Any 2-ply game tree (i.e., each player moves once)
- Max's best move at root:  $a_1$  (leading to highest minimax value of 3)
- ...because Min's best reply will be  $b_1$  (leading to lowest minimax value / utility of 3)





# Minimax (contd.)

## Algorithm and properties

```
function Minimax-Decision(state) returns an action
  inputs: state, current state in game
  return the a in Actions(state) maximizing Min-Value(Result(a, state))
```

```
function Max-Value(state) returns a utility value
  if Terminal-Test(state) then return Utility(state)
  v ← -∞
  for a, s in Successors(state) do v ← Max(v, Min-Value(s))
  return v
```

```
function Min-Value(state) returns a utility value
  if Terminal-Test(state) then return Utility(state)
  v ← ∞
  for a, s in Successors(state) do v ← Min(v, Max-Value(s))
  return v
```

$$\text{MINIMAX}(s) = \begin{cases} \text{UTILITY}(s) & \text{if TERMINAL-TEST}(s) \\ \max_{a \in \text{Actions}(s)} \text{MINIMAX}(\text{RESULT}(s, a)) & \text{if PLAYER}(s) = \text{MAX} \\ \min_{a \in \text{Actions}(s)} \text{MINIMAX}(\text{RESULT}(s, a)) & \text{if PLAYER}(s) = \text{MIN} \end{cases}$$

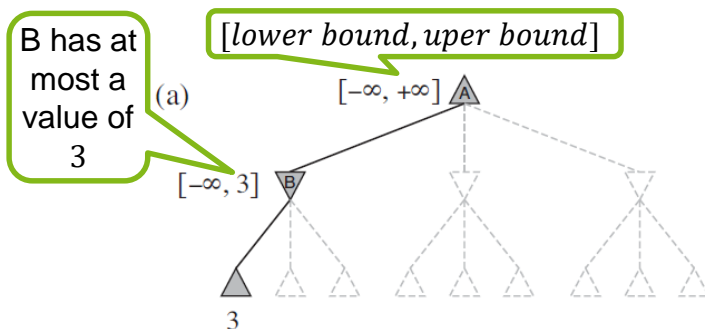
Min-Value
  
Max-Value

## Properties

- Complete? **Yes**, if tree is finite (e.g., chess has specific rules for this)
- Optimal? **Yes**, against an optimal opponent (Otherwise?)
- Time complexity?  $O(b^m)$
- Space complexity?  $O(bm)$  (for depth-first exploration)
- For chess,  $b = 35$ ,  $m = 100$  for “reasonable” games
  - ➔ exact solution completely **infeasible**; but do we need to explore every path?

# $\alpha$ - $\beta$ pruning example

Overcoming exponential ( $b^m$ ) number of states to be explored



Recursion: B

$$\alpha = -\infty$$

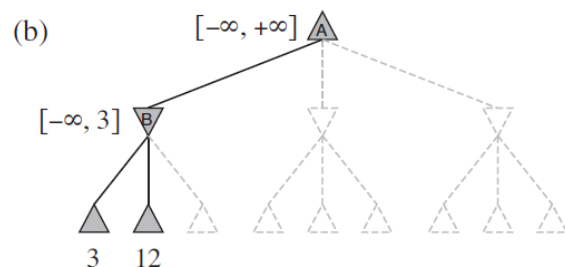
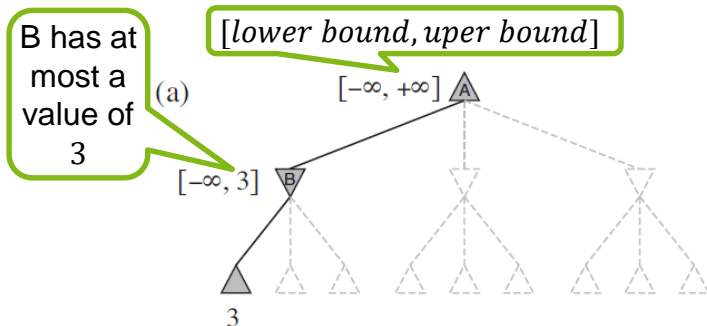
$$\beta = 3$$

Successively tightening bounds on minimax values

- $\alpha$  is the **best value** (to Max) found so far in current subtree of a Max node ( $A \rightarrow \alpha$ )
- If any node  $v$  is worse than  $\alpha$ , Max will not choose it  
 **$\rightarrow$  prune that branch**
- Similarly:  $\beta$  is best score Min is assured of in current subtree of a Min node ( $B, C, D \rightarrow \beta$ )

# $\alpha$ - $\beta$ pruning example

Overcoming exponential ( $b^m$ ) number of states to be explored



Recursion: B

$\alpha = -\infty$

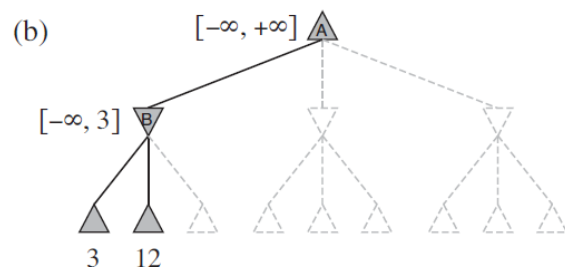
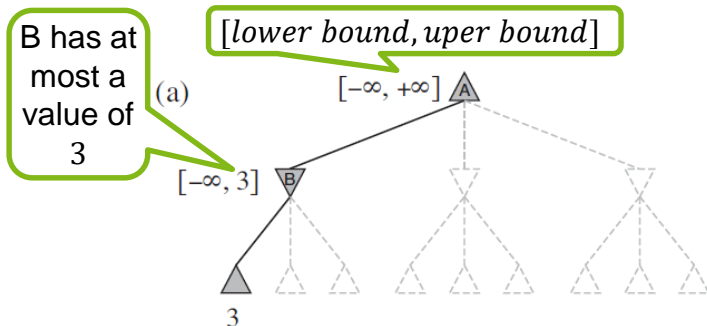
$\beta = 3$

Successively tightening bounds on minimax values

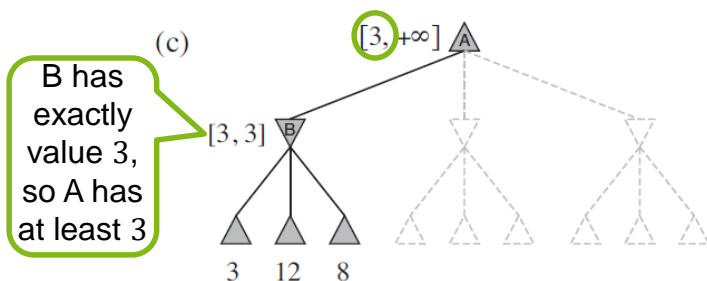
- $\alpha$  is the **best value** (to Max) found so far in current subtree of a Max node ( $A \rightarrow \alpha$ )
- If any node  $v$  is worse than  $\alpha$ , Max will not choose it  
**→ prune that branch**
- Similarly:  $\beta$  is best score Min is assured of in current subtree of a Min node ( $B, C, D \rightarrow \beta$ )

# $\alpha$ - $\beta$ pruning example

Overcoming exponential ( $b^m$ ) number of states to be explored



Recursion: A  
 $\alpha = 3$   
 $\beta = \infty$

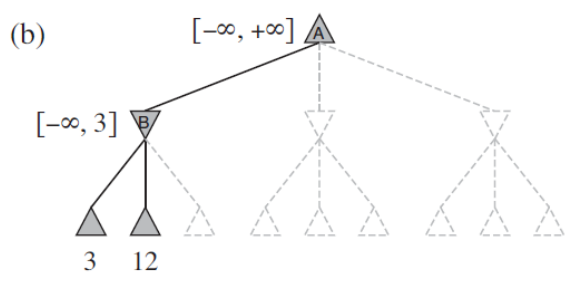
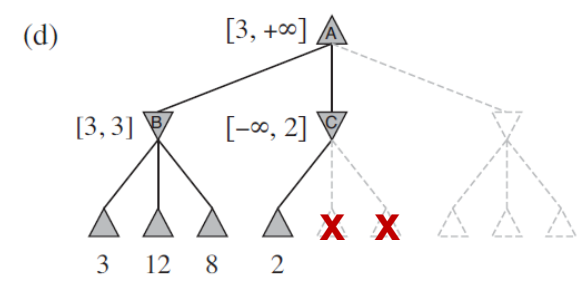
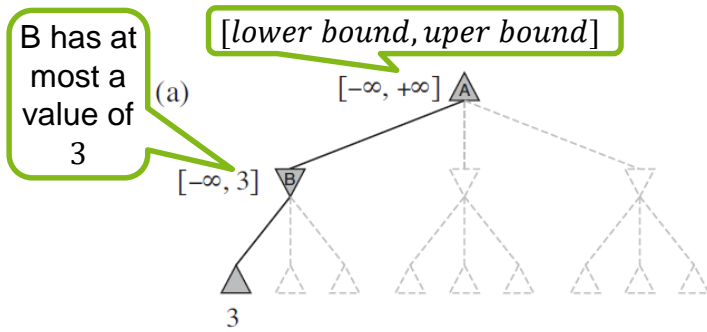


Successively tightening bounds on minimax values

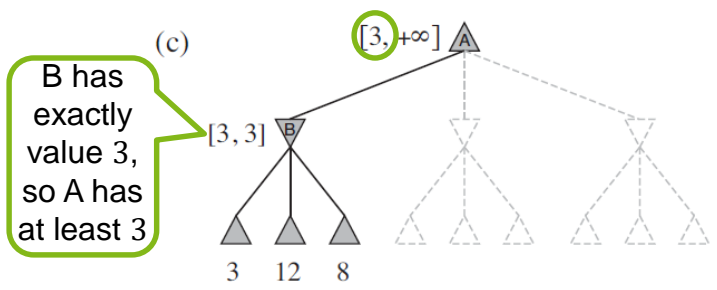
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- If any node  $v$  is worse than  $\alpha$ , Max will not choose it  $\rightarrow$  **prune that branch**
- Similarly:  $\beta$  is best score Min is assured of in current subtree of a Min node ( $B, C, D \rightarrow \beta$ )

# $\alpha$ - $\beta$ pruning example

Overcoming exponential ( $b^m$ ) number of states to be explored



Recursion: C  
 $\alpha = 3$   
 $\beta = 2$

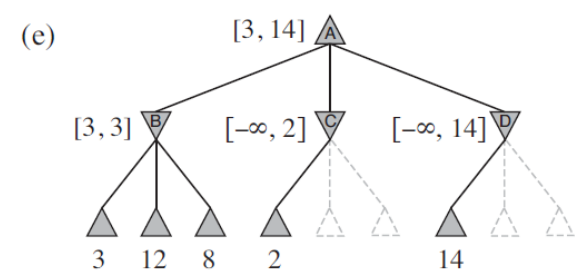
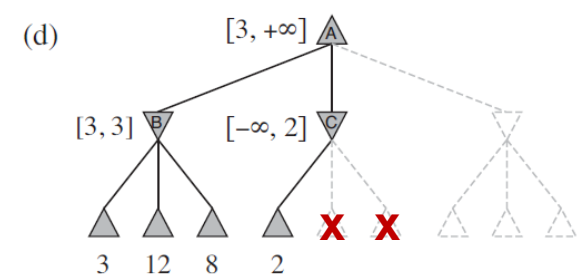
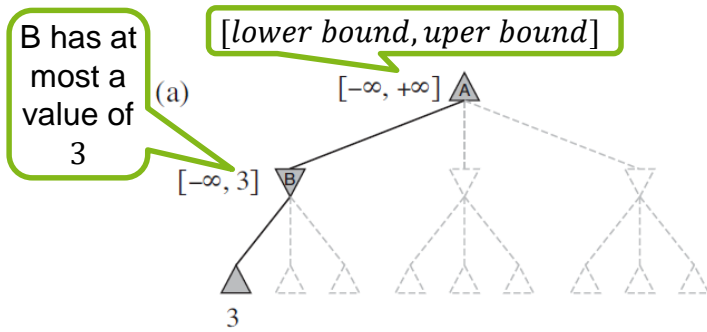


Successively tightening bounds on minimax values

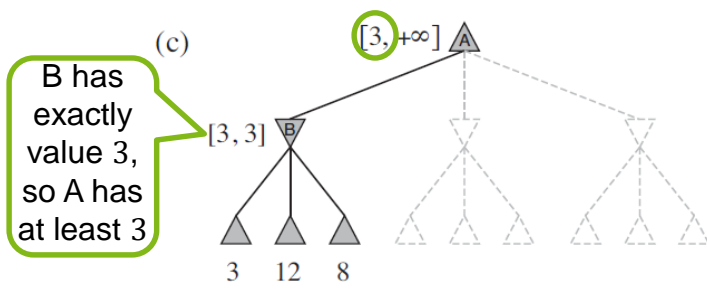
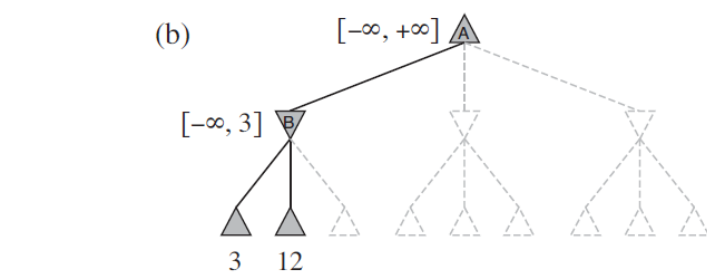
- $\alpha$  is the **best value** (to Max) found so far in current subtree of a Max node ( $A \rightarrow \alpha$ )
- If any node  $v$  is worse than  $\alpha$ , Max will not choose it  $\rightarrow$  **prune that branch**
- Similarly:  $\beta$  is best score Min is assured of in current subtree of a Min node ( $B, C, D \rightarrow \beta$ )

# $\alpha$ - $\beta$ pruning example

Overcoming exponential ( $b^m$ ) number of states to be explored



Recursion: D  
 $\alpha = 3$   
 $\beta = 14$

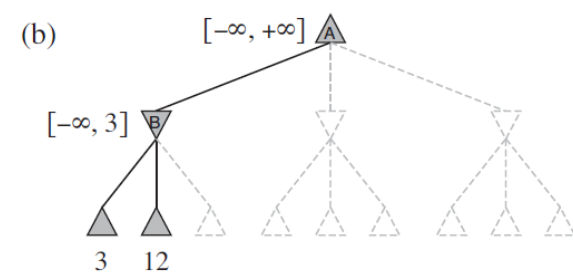
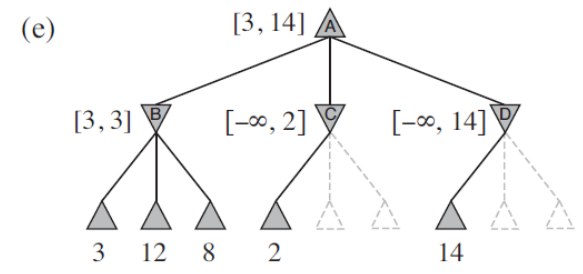
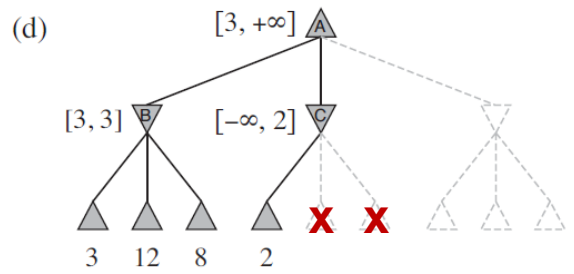
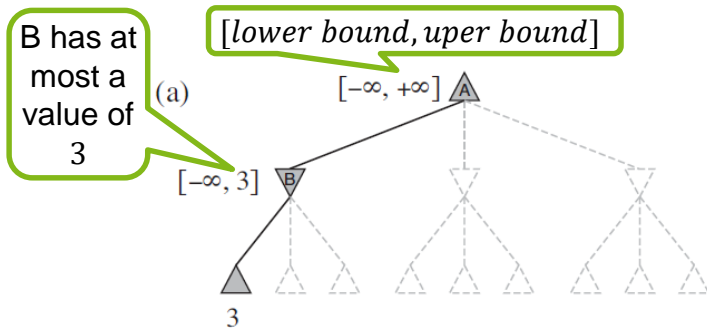


Successively tightening bounds on minimax values

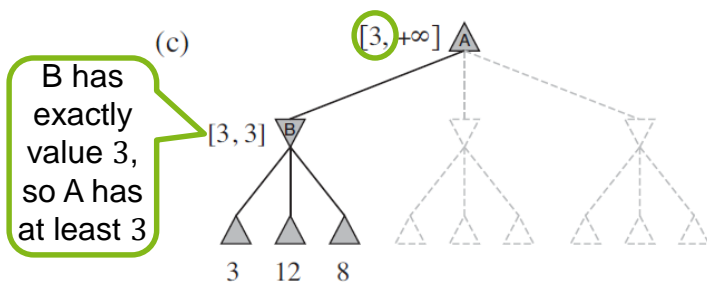
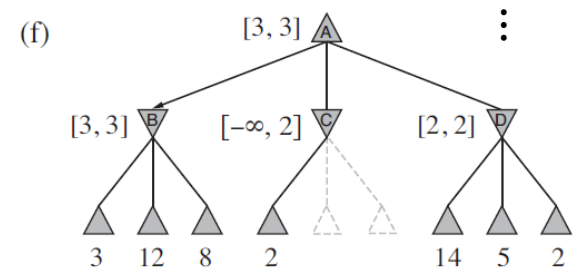
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- Similarly:  $\beta$  is best score Min is assured of in current subtree of a Min node ( $B, C, D \rightarrow \beta$ )

# $\alpha$ - $\beta$ pruning example

Overcoming exponential ( $b^m$ ) number of states to be explored



Recursion: A  
 $\alpha = 3$   
 $\beta = 3$



Successively tightening bounds on minimax values

- $\alpha$  is the **best value** (to Max) found so far in current subtree of a Max node ( $A \rightarrow \alpha$ )
- If any node  $v$  is worse than  $\alpha$ , Max will not choose it  $\rightarrow$  **prune that branch**
- Similarly:  $\beta$  is best score Min is assured of in current subtree of a Min node ( $B, C, D \rightarrow \beta$ )

# $\alpha$ - $\beta$ pruning (contd.)

## Algorithm and properties (changes to minimax in *bold-italic*)

```
function Alpha-Beta-Search(state) returns an action
  v ← Max-Value(state, -∞, ∞)
  return the a in Actions(state) with value v
```

```
function Max-Value(state,  $\alpha$ ,  $\beta$ ) returns a utility value
  if Terminal-Test(state) then return Utility(state)
  v ← -∞
  for a in Actions(state) do
    v ← Max(v, Min-Value(Result(state, a),  $\alpha$ ,  $\beta$ ))
    if v ≥  $\beta$  then return v
     $\alpha$  ← Max( $\alpha$ , v)
  return v
```

```
function Min-Value(state,  $\alpha$ ,  $\beta$ ) returns a utility value
  if Terminal-Test(state) then return Utility(state)
  v ← ∞
  for a in Actions(state) do
    v ← Min(v, Max-Value(Result(state, a),  $\alpha$ ,  $\beta$ ))
    if v ≤  $\alpha$  then return v
     $\beta$  ← Min( $\beta$ , v)
  return v
```

## Properties

- Pruning **does not affect final result**
- Good **move ordering improves** effectiveness of pruning
- With “perfect ordering”, time complexity =  $O(b^{m/2})$ 
  - **doubles solvable depth**
  - a simple example of the value of reasoning about which computations are relevant (**metareasoning**)
  - unfortunately,  $35^{100/2}$  (for chess) is still impossible!



### 3. RESOURCE LIMITS AND OTHER DIFFICULTIES

Adversarial search (contd.)

# Resource limits

## Towards real-world conditions

### Standard approach

- Use **Cutoff-Test** instead of Terminal-Test  
e.g., depth limit (perhaps add **quiescence search**: only cut off search at positions that don't drastically change their value in the near future, e.g. captures in chess; otherwise continue search)
- Use **Eval** instead of Utility  
i.e., evaluation function that estimates desirability of position
- **Lookup** of start/end games

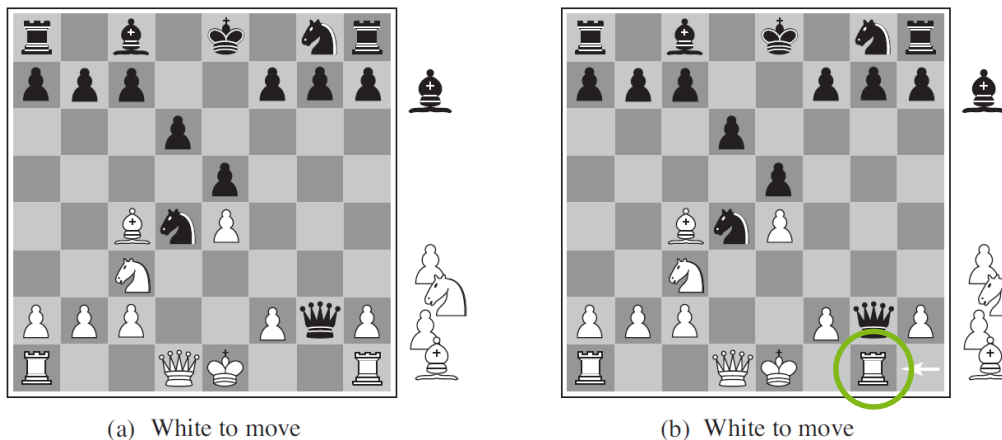
### Example

- Suppose we have 100 seconds, explore  $10^4$  nodes/second  
→  $10^6$  nodes per move  $\approx 35^{8/2}$
- $\alpha$ - $\beta$  reaches depth 8  
→ pretty good chess program



# Eval(uation) functions

## Designing or learning effective cutoff tests



Example: The two chess positions **differ only in the position of the rook** (“Turm”) at lower right.  
 In (a), **Black has an advantage** of a knight (“Springer”) and two pawns (“Bauern”) → should be enough to win.  
 In (b), **White will capture the queen** (“Dame”) → should be strong enough to win.

For chess, typically **linear weighted sum of features**

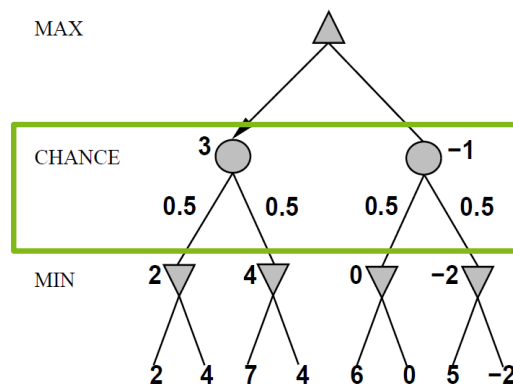
- $Eval(s) = w_1 f_1(s) + w_2 f_2(s) + \dots + w_n f_n(s)$
- Example:  
 $w_1 = 9$  and  $f_1(s) = (\text{number of white queens}) - (\text{number of black queens})$
- **Can be learned** with machine learning techniques

# Nondeterministic (stochastic) games

Chance is introduced by e.g. dice-rolling or card-shuffling

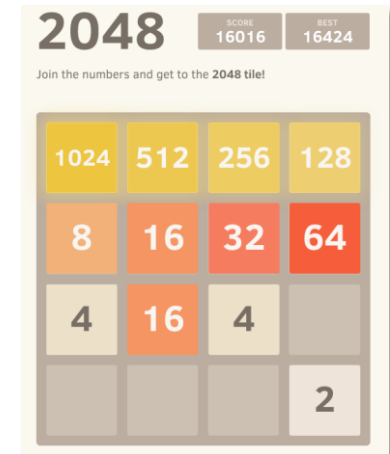
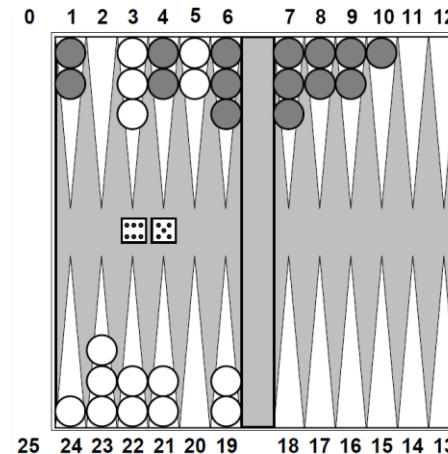
## Simplified example

- A game with coin-flipping
- Nondeterminism is handled by an **additional level in the tree**, consisting of **chance nodes**



## Real-world example

- **2048**: numbers appear with probability  $P(2) = \frac{9}{10}$  and  $P(4) = \frac{1}{10}$  at random free board positions
- **Backgammon**: Before each move, dice-rolls determine the legal moves



# Expectiminimax – maximizing expected value

## Algorithm and properties

```
function ExpectiMinimax-Decision(state) returns an action
  inputs: state, current state in game
  return a in Actions(state) maximizing
    ExpectiMinimax-Value(Result(a, state))

function ExpectiMinimax-Value(state) returns a utility value
  if Terminal-Test(state) then
    return Utility(state)
  if state is a Max node then
    return highest ExpectiMinimax-Value of Successors(state)
  if state is a Min node then
    return lowest ExpectiMinimax-Value of Successors(state)
  if state is a chance node then
    return average of ExpectiMinimax-Value of Successors(state)
```

$$\text{EXPECTIMINIMAX}(s) = \begin{cases} \text{UTILITY}(s) & \text{if } \text{TERMINAL-TEST}(s) \\ \max_a \text{EXPECTIMINIMAX}(\text{RESULT}(s, a)) & \text{if } \text{PLAYER}(s) = \text{MAX} \\ \min_a \text{EXPECTIMINIMAX}(\text{RESULT}(s, a)) & \text{if } \text{PLAYER}(s) = \text{MIN} \\ \sum_r P(r) \text{EXPECTIMINIMAX}(\text{RESULT}(s, r)) & \text{if } \text{PLAYER}(s) = \text{CHANCE} \end{cases}$$

## Properties

- Algorithm **works** just like **Minimax** – except chance-nodes are also handled
- Expectiminimax gives **perfect play**
- In case of only **1 player**, Expectiminimax becomes **Expectimax**
- Time complexity:  $O(b^m n^m)$  (where  $n$  is the number of distinct random events, e.g. dice rolls)
  - **Possibilities** are **multiplied enormously** in games of chance
  - Simultaneously, **no** likely sequences exist to do effective  **$\alpha$ - $\beta$  pruning**

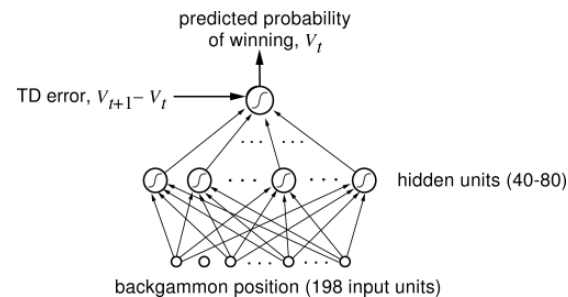
# Nondeterministic games in practice

## Example Backgammon

- Dice rolls increase  $b$  (21 possible rolls with 2 dice)
- Ca. 20 legal moves (can be 6,000 with 1-1 roll)
  - at depth 4:  $20 \times (21 \times 20)^3 \approx 1.2 \times 10^9$  nodes
- As depth increases, probability of reaching a given node shrinks
  - value of lookahead is diminished (see next slide for consequences)

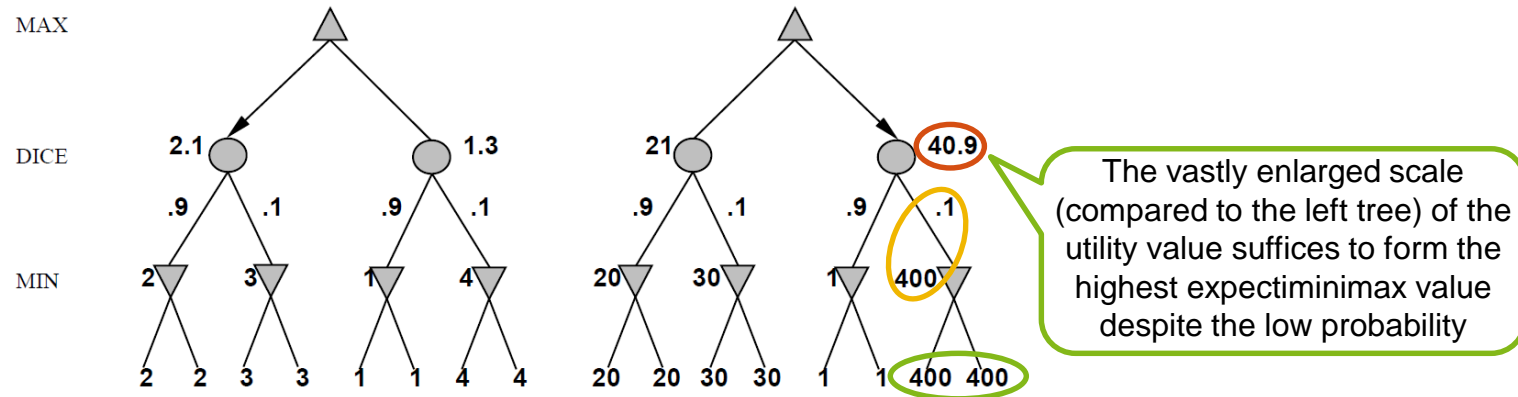
## But

- «TDGammon» (Tesauro, 1992) uses depth-2 search  $\approx$  world-champion level
  - uses neural network and **reinforcement learning** (RL) to train Eval function via games against itself



- → see also <http://webdocs.cs.ualberta.ca/~sutton/book/ebook/node108.html>

# Digression: Exact Eval values do matter ...for nondeterministic games



- Behavior is **preserved only by positive linear transformation** of Eval
- Hence **Eval** should be **proportional to the expected payoff**

→ Exact values **don't matter for deterministic** games → see appendix

# Games of imperfect information

E.g., card games, where opponent's initial cards are unknown

First idea («[averaging over clairvoyance](#)»)

- Typically we can **calculate a probability** for **each possible deal**
  - Seems just like having one big dice roll at the beginning of the game \*
  - **Idea:** compute the Minimax value of each action in each deal;  
then, choose the action with highest expected value over all deals (i.e., **Expectiminimax**) \*

**Monte Carlo** simulation alone does not always suffice: it **handles randomness well** (e.g. in the deal), but **not strategy** that looks like randomness (e.g. an adversary's move in Kriegspiel)

«GIB» - Ginsberg's Intelligent **Bridge**player

- Current best bridge program (Ginsberg, 1999)
- Approximates idea above by [Monte Carlo simulation](#) of some of the possible deals to be computationally feasible
  1. **Generate 100 deals** consistent with bidding information
  2. **Pick the action that wins most tricks** on average



Skat: a wonderful partially observable stochastic card game with some similarity to Bridge

→ see <https://www.pagat.com/schafk/skat.html>



## \* ) Proper analysis of averaging over clairvoyance

Intuition might mislead

- That the value of an action is the **average of its values in all possible states** is **WRONG** (→ see appendix)
- With partial observability, value of an action depends on the **belief state** the agent will end up in

Maintaining its belief state is a **core function of any intelligent system** in partially observable (~real world) environments → see AIMA, ch. 4.4

The agent's **current information about all the possible physical states** it might be in (given the sequence of actions and percepts up to that point)

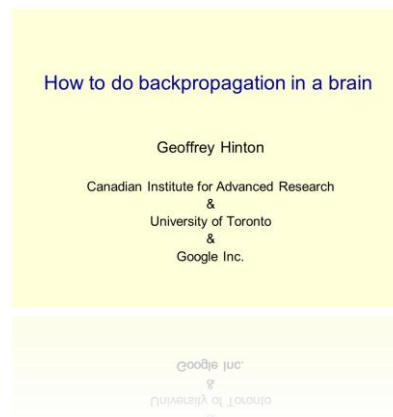
Dealing with belief states

- Generate and search a tree of belief states
- Leads to rational behaviors such as
  - Acting to obtain information
  - Signaling to one's partner
  - Acting randomly to minimize information disclosure

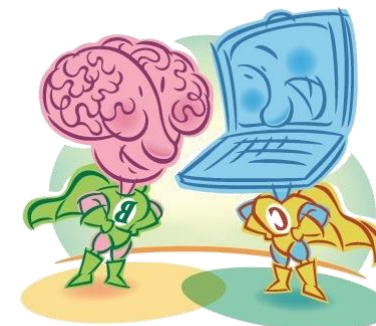
# Where's the intelligence?

## Man vs. machine

- Searching over the right state spaces leads to **emergent “clever” behavior** in games (information gathering, alliances, bluffs)
- Many local and adversarial search methods can be **enhanced by learning** (e.g., learn good Eval functions by playing games against oneself)
- The brain might perform local (hill climbing) search as well  
→ see <https://www.cs.toronto.edu/~hinton/backpropincortex2014.pdf>



- Hill-climbing, minimax,  $\alpha$ - $\beta$ , etc. are still pure computation
- Intelligent behavior emerges from their composition on **suitable** data structures



# Review

- **Local search** algorithms **evolve** a **small number of states** towards better utility (typical: 1 state)
- In **continuous space**, local search by **linear programming** or **convex optimization** is extremely efficient in practice (polynomial time complexity!)
- In non-deterministic environments, **keeping track of one's belief state** is paramount
- Games are fun to work on – and dangerous
- They illustrate several important points about AI
  - **perfection** is **unattainable** → must **approximate**
  - good idea to **think about what to think** about (metareasoning → pruning)
  - **uncertainty** constrains the assignment of values to states
  - optimal decisions depend on **information state** (belief state), not real state
- Games are to AI as grand prix racing is to automobile design





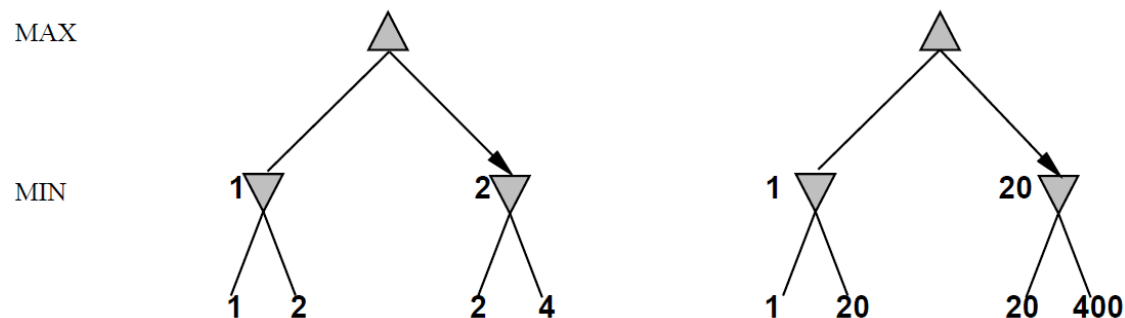
# APPENDIX

# Pseudocode for hill climbing search and simulated annealing

```
function Hill-Climbing(problem) returns a state that is a local maximum
  inputs: problem, a problem
  local variables: current, a node
                   neighbour, a node
  current ← Make-Node(Initial-State[problem])
  loop do
    neighbour ← a highest-valued successor of current
    if Value[neighbour] ≤ Value[current] then return State[current]
    current ← neighbour
  end
```

```
function Simulated-Annealing(problem, schedule) returns a solution state
  inputs: problem, a problem
         schedule, a mapping from time to "temperature"
  local variables: current, a node
                 next, a node
                 T, a "temperature" controlling the probability of downward steps
  current ← Make-Node(Initial-State[problem])
  for t ← 1 to ∞ do
    T ← schedule[t]
    if T = 0 then return current
    next ← a randomly selected successor of current
    ΔE ← Value[next]-Value[current]
    if ΔE > 0 then current ← next
    else current ← next only with probability  $e^{-\frac{\Delta E}{T}}$ 
```

# Digression: Exact Eval values don't matter ...for deterministic games



- Behavior is preserved under any **monotonic** transformation of Eval
- Only the **order matters**: payoff in deterministic games acts as an **ordinal utility** function

# Deterministic games in practice

## Checkers

- «*Chinook*» ended 40-year-reign of human world champion Marion Tinsley in 1994
- Used an endgame database defining perfect play for all positions involving 8 or fewer pieces on the board (total: 443,748,401,247)



## Chess

- «*Deep Blue*» defeated human world champion Gary Kasparov in a six-game match in 1997
- Searches 200 million positions per second, uses very sophisticated evaluation, and undisclosed methods for extending some lines of search up to 40 ply



## Othello

- Human champions refuse to compete against computers, who are **too good**



## Go

- 2010: human champions refuse to compete against computers, who are **too bad**
- 2016: Google DeepMind's «*AlphaGo*» unexpectedly defeats world champion Lee Sedol
- $b > 300 \rightarrow$  most programs use pattern knowledge bases to suggest plausible moves



# Expectations & expected values

adapted from U Washington`s CSE473, lecture 8

The expectation operator  $E()$

- We can define a function  $f(X)$  of a **random variable**  $X$
- The **expected value** of a function is its **average** value **under the probability distribution over the function's inputs**:

$$E(f(X)) = \sum_x f(X = x)P(X = x)$$

## Example

- How long to drive to the airport?
- Driving time  $D$  (in mins) as a function of traffic  $T$ :  
 $D(T = none) = 20, D(T = light) = 30, D(T = heavy) = 60$
- What is your expected driving time?  
Let probability  $P(T) = \{none: 0.25, light: 0.5, heavy: 0.25\}$   
→  $E(D(T)) = D(none)P(none) + D(light)P(light) + D(heavy)P(heavy)$   
→  $E(D(T)) = 20 * 0.25 + 30 * 0.5 + 60 * 0.25 = 35$  mins



# Why averaging over clairvoyance is wrong

## A common sense example of a journey

### Day 1

- Road A leads to a small heap of gold pieces
  - Road B leads to a fork:
    - the left fork **leads to** a bigger heap of gold;
    - take the right fork and you'll be run over by a bus.
- „B“ (and then „left“) is the optimal choice

### Day 2

- Road A leads to a small heap of gold pieces
  - Road B leads to a fork:
    - **take** the left fork and you'll be run over by a bus;
    - the right fork leads to a bigger heap of gold.
- „B“ (and then „right“) is the optimal choice

### Day 3

- Road A leads to a small heap of gold pieces
  - Road B leads to a fork:
    - **guess** correctly and you'll find a bigger heap of gold;
    - guess incorrectly and you'll be run over by a bus.
- „B“ still seems optimal; but this **ignores** the resulting **belief state** (that includes ignorance & possibility of death!)



Averaging over clairvoyance **will never** select actions to **gather information** because it assumes future states to be of perfect knowledge after the initial deal.